

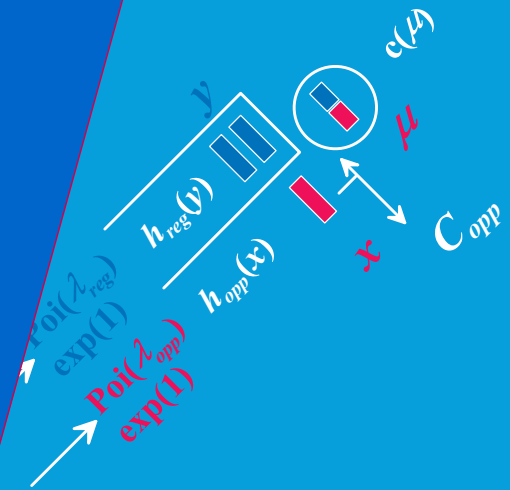
Head-of-line processor sharing: optimal control

Sandra van Wijk

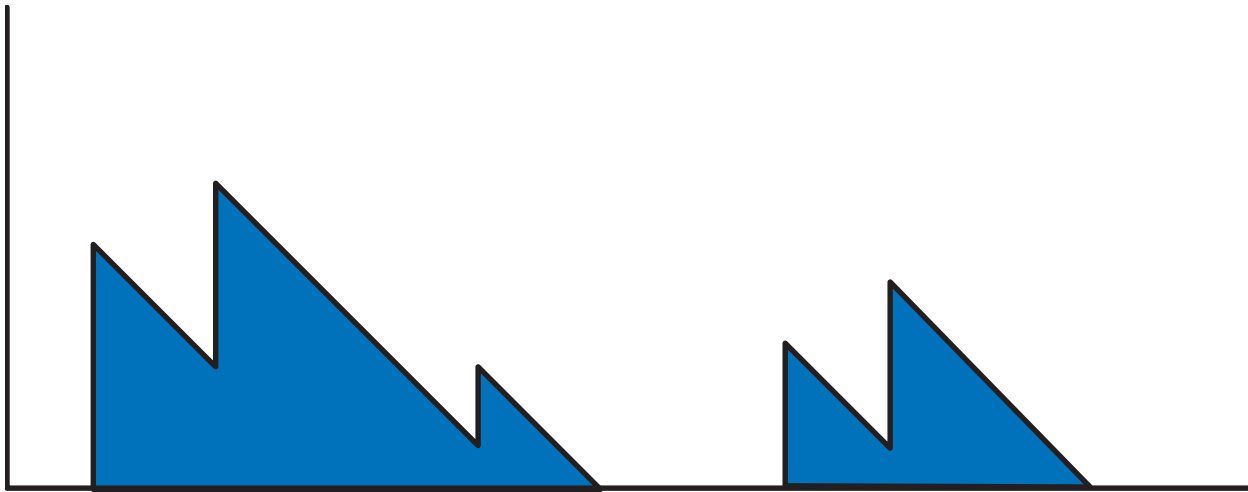
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Joint work with Çağdas Büyükkaramikli

Advisors: Ivo Adan and Geert-Jan van Houtum

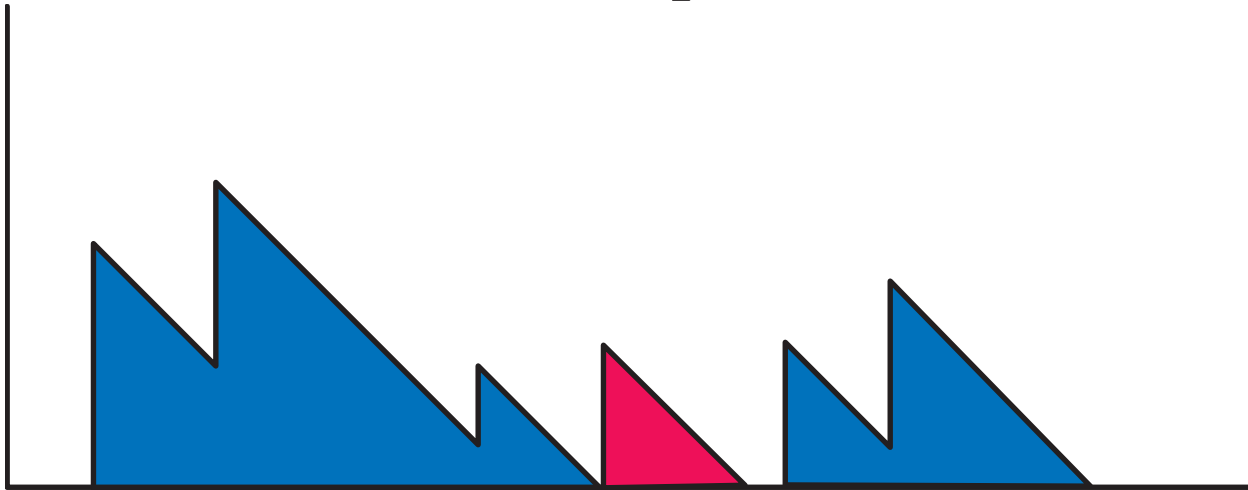


Workload in $M/M/1$ queue



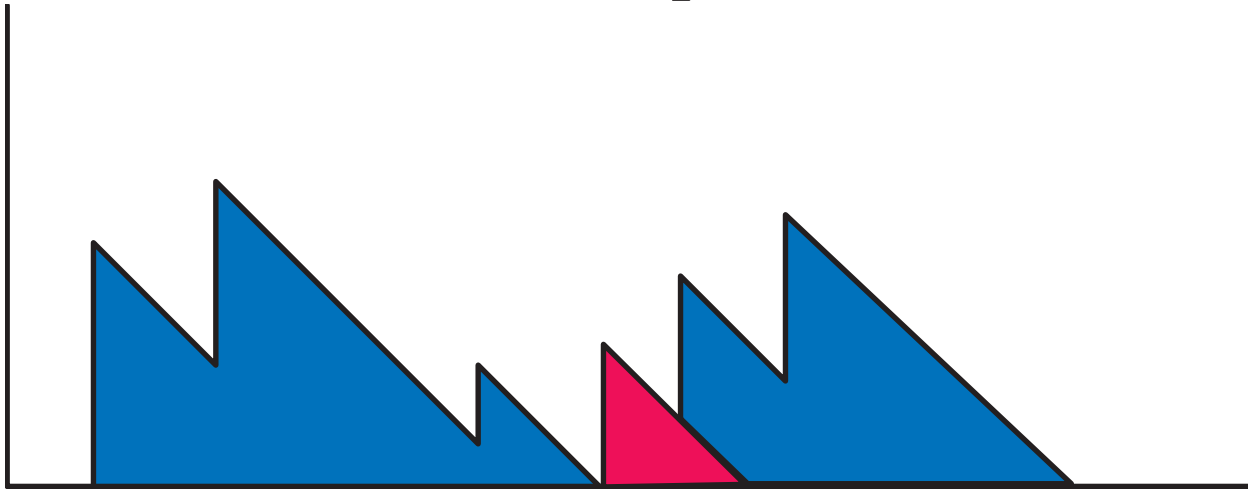
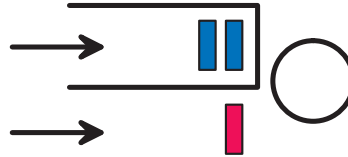
Application: e.g. production system

Workload in $M/M/1$ queue



Extra revenue for serving opportunity customers

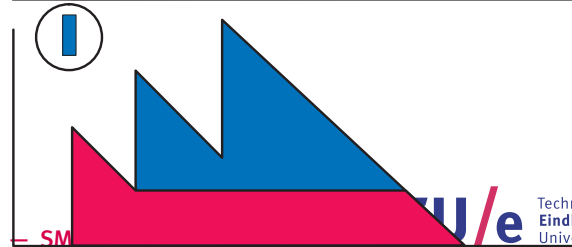
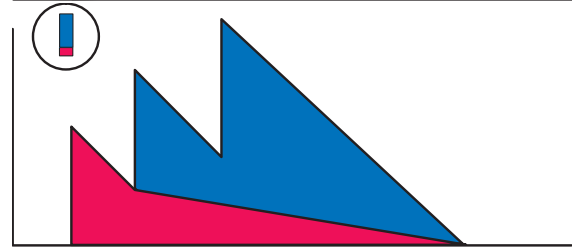
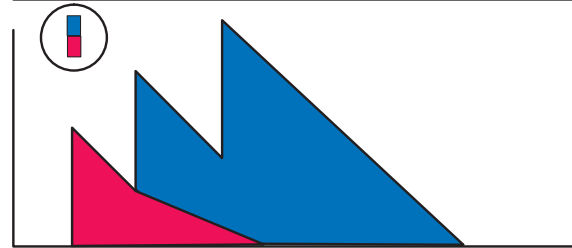
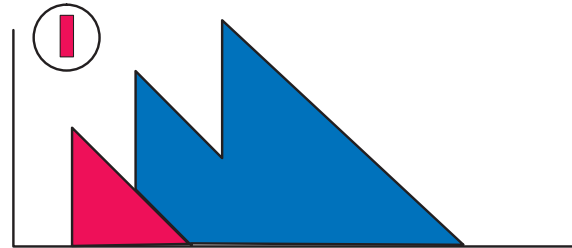
Workload in $M/M/1$ queue

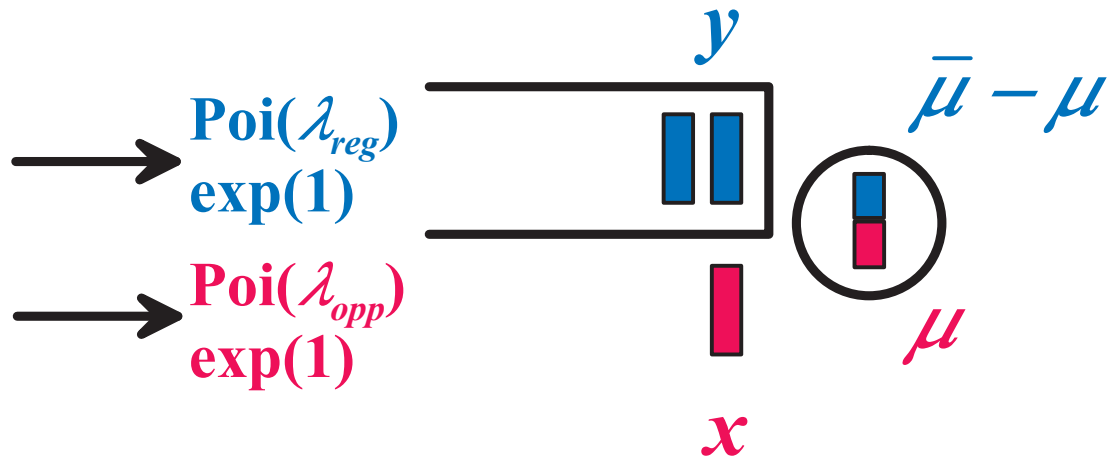


However, waiting times (and hence costs) of regular customers increase

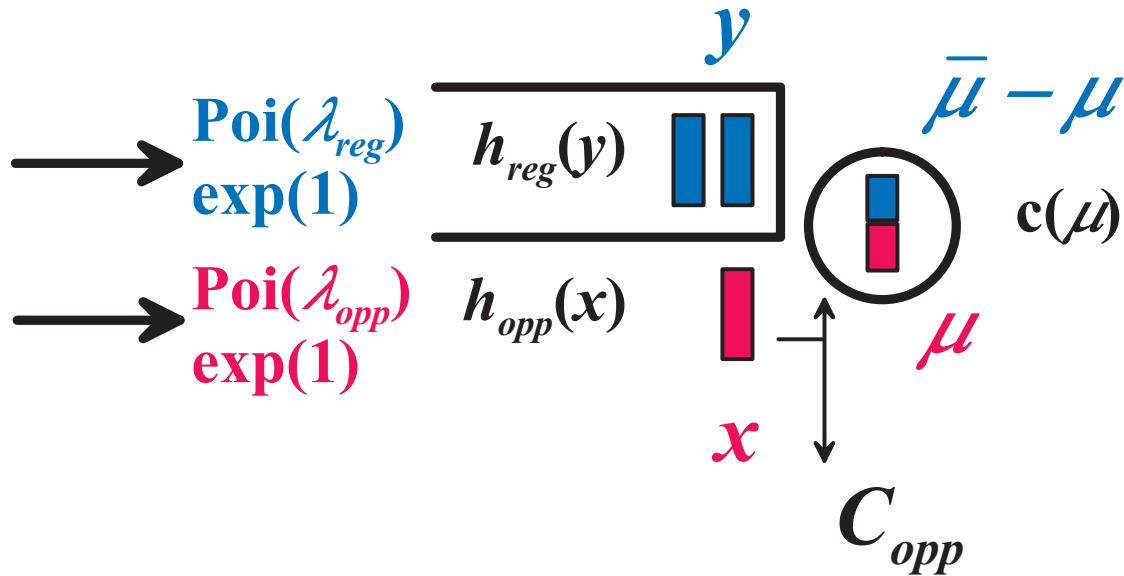
Workload in $M/M/1$ queue

- Regular and opportunity customers
- Head-of-line processor sharing





Research Question: What is optimal control policy?



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Markov Decision Problem

Value function:

$$\begin{aligned} V_{n+1}(x, y) &= h_{opp}(x) + h_{reg}(y) + \frac{1}{\lambda_{opp} + \lambda_{reg} + \bar{\mu} + \alpha} \\ &\quad \left(\lambda_{reg} V_n(x, y + 1) \right. \\ &\quad \left. + \lambda_{opp} \begin{cases} \min\{V_n(x + 1, y), V_n(x, y) + C_{opp}\} & \text{if } x = 0 \\ V_n(x, y) + C_{opp} & \text{if } x = 1 \end{cases} \right. \\ &\quad \left. + \min_{\mu \in [0, \bar{\mu}]} \left\{ c'(\mu) + \mu V_n((x - 1)^+, y) + (\bar{\mu} - \mu) V_n(x, (y - 1)^+) \right\} \right), \\ V_0 &\equiv 0. \end{aligned}$$

Structural properties of value function imply optimal control policy structure.

Value function

Structural property: Multimodularity (MM, cf. Hajek, 1985)

Multimodularity (for 2 dimensions):

$$\text{Supermodularity: } f(x, y) + f(x + 1, y + 1) \geq f(x + 1, y) + f(x, y + 1),$$

$$\text{Superconvexity(1,2): } f(x + 2, y) + f(x, y + 1) \geq f(x + 1, y) + f(x + 1, y + 1),$$

$$\text{Superconvexity(2,1): } f(x, y + 2) + f(x + 1, y) \geq f(x, y + 1) + f(x + 1, y + 1),$$

which implies:

$$\text{Convexity(1): } f(x, y) + f(x + 2, y) \geq 2f(x + 1, y),$$

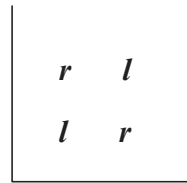
$$\text{Convexity(2): } f(x, y) + f(x, y + 2) \geq 2f(x, y + 1).$$

Value function

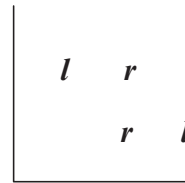
Structural property: Multimodularity (MM, cf. Hajek, 1985)

Multimodularity (for 2 dimensions):

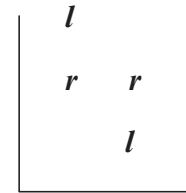
$$l \geq r$$



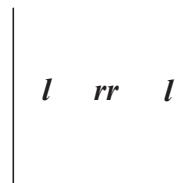
Supermodularity



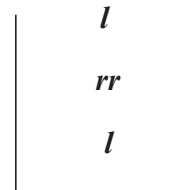
Superconvexity(1,2)



Superconvexity(2,1)



Convexity(1)



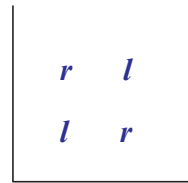
Convexity(2)

Value function

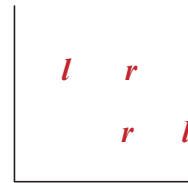
Structural property: Multimodularity (MM, cf. Hajek, 1985)

Multimodularity (for 2 dimensions):

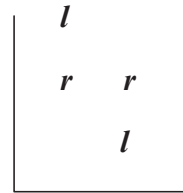
$$l \geq r$$



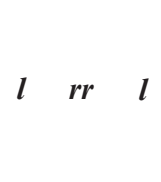
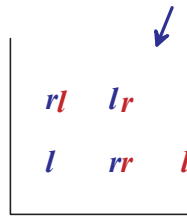
Supermodularity



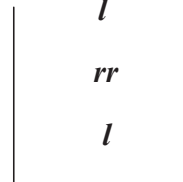
Superconvexity(1,2)



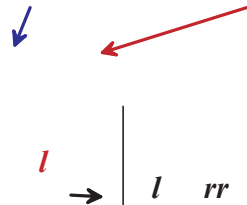
Superconvexity(2,1)



Convexity(1)



Convexity(2)



Value function

Prove that V_n is MM by induction on n :

- $V_0 \equiv 0$ is MM,
- Assume V_n is MM, show that V_{n+1} is MM.

Recall:

$$\begin{aligned} V_{n+1}(x, y) &= h_{opp}(x) + h_{reg}(y) + \frac{1}{\lambda_{opp} + \lambda_{reg} + \bar{\mu} + \alpha} \\ &\quad \left(\lambda_{reg} V_n(x, y + 1) \right. \\ &\quad \left. + \lambda_{opp} \begin{cases} \min\{V_n(x + 1, y), V_n(x, y) + C_{opp}\} & \text{if } x = 0 \\ V_n(x, y) + C_{opp} & \text{if } x = 1 \end{cases} \right) \\ &\quad + \min_{\mu \in [0, \bar{\mu}]} \left\{ c'(\mu) + \mu V_n((x - 1)^+, y) + (\bar{\mu} - \mu) V_n(x, (y - 1)^+) \right\}, \end{aligned}$$

Value function

Rewrite value function using event operators (cf. Koole, 2006)

$$\begin{aligned} V_{n+1}(x, y) \\ = T_{\text{costs}} \left(T_{\text{unif}} \left(T_{CA(1)} V_n(x, y), T_{A(2)} V_n(x, y), \tilde{T}_{CTD(1)} V_n(x, y) \right) \right) \end{aligned}$$

Event operators:

- Arrivals of opportunity customers (decision!)

$$T_{CA(1)} f(x, y) = \min\{V_n(x + 1, y), V_n(x, y) + C_{opp}\}$$

- Arrivals of regular customers

$$T_{A(2)} f(x, y) = V_n(x, y + 1)$$

Value function

Rewrite value function using event operators (cf. Koole, 2006)

$$\begin{aligned} V_{n+1}(x, y) \\ &= T_{\text{costs}} \left(T_{\text{unif}} \left(T_{CA(1)} V_n(x, y), T_{A(2)} V_n(x, y), \tilde{T}_{CTD(1)} V_n(x, y) \right) \right) \end{aligned}$$

Event operators (ctd.):

- Service completions (decision!)

$$\tilde{T}_{CTD(1)} f(x, y) = \min_{\mu \in [0,1]} \left\{ c(\mu) + \mu V_n((x-1)^+, y) + (1-\mu) V_n(x, (y-1)^+) \right\}$$

Value function

Rewrite value function using event operators (cf. Koole, 2006)

$$\begin{aligned} V_{n+1}(x, y) \\ = T_{\text{costs}} \left(T_{\text{unif}} \left(T_{CA(1)} V_n(x, y), T_{A(2)} V_n(x, y), \tilde{T}_{CTD(1)} V_n(x, y) \right) \right) \end{aligned}$$

Event operators (ctd.):

- Costs

$$T_{\text{costs}} f(x, y) = h_{\text{opp}}(x) + h_{\text{reg}}(y) + f(x, y)$$

- Uniformization

$$T_{\text{unif}}(f_1, f_2, f_3)(x, y) = \frac{\lambda_{\text{opp}} f_1(x, y) + \lambda_{\text{reg}} f_2(x, y) + \bar{\mu} f_3(x, y)}{\lambda_{\text{opp}} + \lambda_{\text{reg}} + \bar{\mu} + \alpha}$$

Value function

Use known results for operators (Koole, 2006):

- V_n is MM $\Rightarrow T_{CA(1)} V_n$ is MM
- V_n is MM $\Rightarrow T_{A(2)} V_n$ is MM
- V_n is MM $\Rightarrow T_{\text{unif}} V_n$ is MM
- V_n is MM $\Rightarrow T_{\text{costs}} V_n$ is MM

And $\tilde{T}_{CTD(1)}$?

Value function

Compare $\tilde{T}_{CTD(1)}$ to departure operator in **tandem queue**:

$$\tilde{T}_{CTD(1)}f(x, y) = \min_{\mu \in [0, 1]} \left\{ c(\mu) + \mu V_n((x-1)^+, y) + (1-\mu)V_n(x, (y-1)^+) \right\}$$

$$T_{CTD(1)}f(x, y) = \min_{\mu \in [0, 1]} \left\{ c(\mu) + \mu V_n((x-1)^+, y+1) + (1-\mu)V_n(x, y) \right\}$$

Known result: V_n is MM $\Rightarrow T_{CTD(1)}V_n$ is MM

- **Transformation:** $y \rightarrow y - 1$
- check $x = 0, y = 0$ and $x > 0, y = 0$

Then:

- V_n is MM $\Rightarrow \tilde{T}_{CTD(1)}V_n$ is MM

Main result

Value function is **MM**. Implies **optimal policy structure**.

Optimal policy structure for a head-of-line processor sharing model, with adjustable weights and two types of customers:

- **Threshold** T for admitting opportunity customer: accept if $y \leq T$, reject it otherwise.
- The optimal server speed dedicated to the opportunity customer is a **monotone decreasing** function in x .

Example 1: $c(\mu) \equiv 0$

$\lambda_{reg} = 3$, $\lambda_{opp} = 1$, $C_{opp} = 8$, $\bar{\mu} = 10$, $h_{opp}(x) = x$ and $h_{reg}(y) = 0.05 y^2$ if $y < 20$; $h_{reg}(y) = 100 y$ otherwise.

Optimal policy accept/reject opportunity customer:

$x \backslash y$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	1	1	1	1	1	1	1	-	-	-	-	-	-	-	-
1	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-

Threshold policy with $T = 6$.

The optimal fraction $\mu \in [0, 1]$ of server speed for opp. customer:

$x \backslash y$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	1	1	1	1	1	1	1	1	1	1	1	0	0	0	0

So, opportunity customers either gets full attention of the server, or no attention at all, with threshold $y = 10$.

Example 1: $c(\mu) \equiv 0$

- If $c(\mu) \equiv 0$ (or: constant) also threshold policy for optimal server speed for opp. customer.
- Because $\tilde{T}_{CTD(1)}$ minimizes a linear function:

$$\begin{aligned} & \tilde{T}_{CTD(1)} f(x, y) \\ &= \min_{\mu \in [0, 1]} \left\{ c(\mu) + \mu V_n((x - 1)^+, y) \right. \\ & \quad \left. + (1 - \mu) V_n(x, (y - 1)^+) \right\} \end{aligned}$$

Example 2: $c(\mu) \neq 0$

Same parameters, now:

$$c(\mu) = \begin{cases} 0 & \text{if } 0 \leq \mu < 0.25; \\ 0.5 & \text{if } 0.25 \leq \mu < 0.50; \\ 1 & \text{if } 0.50 \leq \mu < 0.75; \\ 1.5 & \text{if } 0.75 \leq \mu \leq 1. \end{cases}$$

Threshold $T = 5$ for accepting opportunity customers.

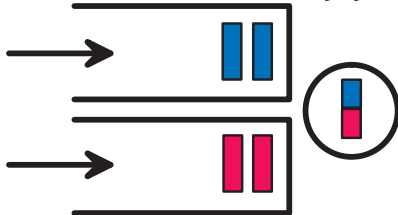
The optimal fraction $\mu \in [0, 1]$ of the service speed for the opp. customer:

$x \backslash y$	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	1	1	1	1	1	1	1	0.75	0.25	0.25	0.25	0.25	0	0	0

Indeed monotone decreasing.

Model extensions

- Accept or reject regular customers
- Queueing of opportunity customers



Further research

- Steady state probability distribution
- Total service rate increases or decreases when the server divides its attention to two customers
- Multiple types of opportunity customers



Conclusion

Optimal policy structure for a head-of-line processor sharing model, with adjustable weights and two types of customers.